

Unifying Cubical Models of Univalent Type Theory

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Univalent Type Theory

▣ Dependent type theory

$$\begin{array}{ll} \Gamma \vdash A \text{ type} & \Gamma \vdash M : A \\ \Gamma \vdash A = B \text{ type} & \Gamma \vdash M = N : A \end{array}$$

Univalent Type Theory

▣ Dependent type theory

$\Gamma \vdash A$ type	$\Gamma \vdash M : A$
$\Gamma \vdash A = B$ type	$\Gamma \vdash M = N : A$

$\Gamma \vdash (a:A) \rightarrow B$ type	function/implication/ \forall
$\Gamma \vdash (a:A) \times B$ type	product/ \exists
$\Gamma \vdash \mathbb{N}$ type	inductive types
$\Gamma \vdash \text{Id}_A(M, N)$ type	equality
$\Gamma \vdash \mathcal{U}$ type	universe(s) of types

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e.g. $\cdot \vdash M : (n : \mathbb{N}) \rightarrow (m : \mathbb{N}) \times \text{Id}_{\mathbb{N}}(m, n)$

Univalent Type Theory

▣ Identity $\text{Id}_A(a, b)$

$$\frac{a : A}{\text{refl}_A a : \text{Id}_A(a, a)} + \frac{a : A \vdash d : P(a, a, \text{refl}_A a)}{\text{elim}(a.d, a_0, a_1, t) : P(a_0, a_1, t)}$$

- ◆ Least reflexive relation (\Rightarrow symmetric, transitive, etc.)
- ◆ “Underdetermined”

$$\text{Id}_{(a:A) \rightarrow B}(f, g) \stackrel{?}{\simeq} (a:A) \rightarrow \text{Id}_B(fa, ga)$$

$$\text{Id}_{\mathcal{U}} A B \simeq ?$$

Univalent Type Theory

Univalence Axiom (Voevodsky)

$$\text{Id}_{\mathcal{U}}(A, B) \simeq (A \simeq B)$$

Equivalence $A \simeq B$

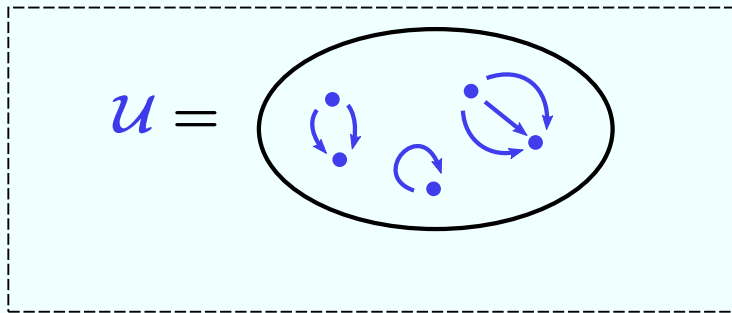
◆ map $f : A \rightarrow B$ with a left and right inverse

$$\begin{array}{ccccccc} & & \text{id} & & & & \\ & \text{---} & \text{---} & \text{---} & \text{---} & \text{---} & \\ & & \beta \parallel & & & & \\ B & \xrightarrow{s} & A & \xrightarrow{f} & B & \xrightarrow{r} & A \\ & & & & \alpha \parallel & & \\ & & & & \text{id} & & \end{array}$$

$$(f, s, \beta, r, \alpha) : A \simeq B$$

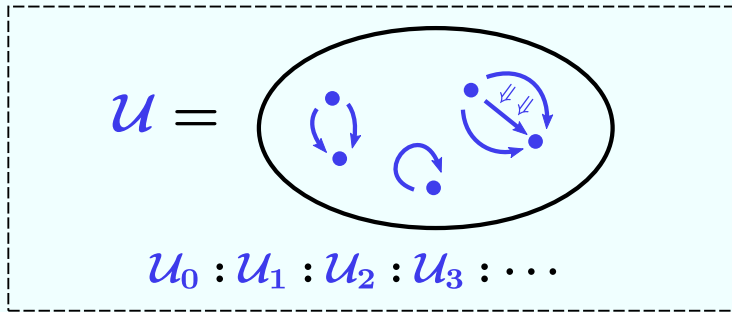
Univalent Type Theory

▣ Identities are not unique



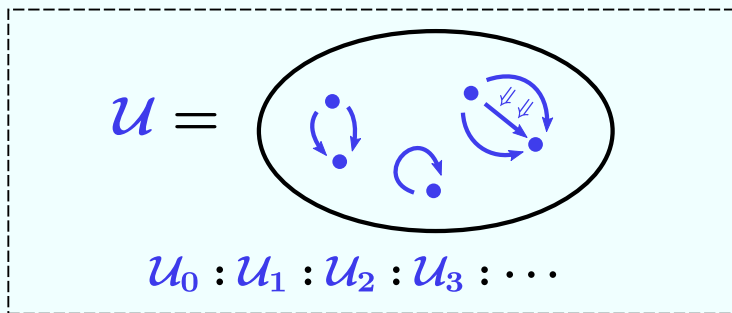
Univalent Type Theory

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Univalent Type Theory

- ▣ Identities are not unique



- ▣ More: add *higher inductive types*
 - ◆ Quotients for proof-relevant identity
 - ◆ Language for *synthetic homotopy theory*

Models of Univalent Type Theory

- ▣ **Simplicial set model**
(Kapulkin & Lumsdaine '12/'18, after Voevodsky)
 - ◆ Classical setting for homotopy theory
 - ◆ Essentially non-constructive
(Bezem, Coquand, & Parmann '15)
- ▣ **Cubical set model**
(Bezem, Coquand, & Huber '13)
 - ◆ First constructive model of univalence
 - ◆ Problems with higher inductive types resolved in Cohen, Coquand, Huber, & Mörtberg '15 and Angiuli, Favonia, & Harper '18 models

Cubical Set Models

▣ Interpret contexts as **cubical sets**

- ◆ family of sets indexed by *interval variable* contexts

$$\Gamma \text{ ctx} \rightsquigarrow \llbracket \Gamma \rrbracket (i_1 : \mathbb{I}, \dots, i_n : \mathbb{I}) \text{ for each } n$$

“maps $[0, 1]^n \rightarrow \Gamma$ ”

$$\llbracket \Gamma \rrbracket (\cdot) = \left\{ \circ \right\} \quad \llbracket \Gamma \rrbracket (i : \mathbb{I}) = \left\{ \begin{array}{c} i \rightarrow \\ \circ \text{---} \circ \end{array} \right\}$$

$$\llbracket \Gamma \rrbracket (i : \mathbb{I}, j : \mathbb{I}) = \left\{ \begin{array}{c} i \rightarrow \\ j \downarrow \\ \text{shaded square} \end{array} \right\}$$

Cubical Set Models

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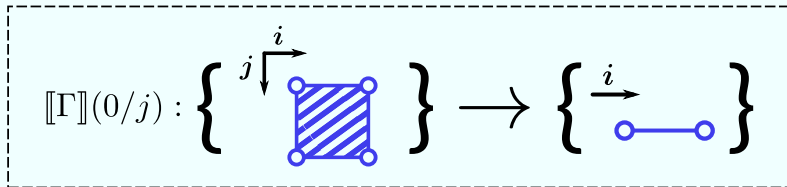
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- ◆ for every $f : (i_1 : \mathbb{I}, \dots, i_n : \mathbb{I}) \rightarrow (j_1 : \mathbb{I}, \dots, j_m : \mathbb{I})$
in some fixed class of interval maps,

$$\llbracket \Gamma \rrbracket (f) : \llbracket \Gamma \rrbracket (j_1 : \mathbb{I}, \dots, j_m : \mathbb{I}) \rightarrow \llbracket \Gamma \rrbracket (i_1 : \mathbb{I}, \dots, i_n : \mathbb{I})$$



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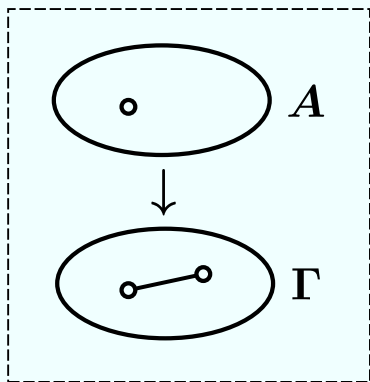
▣ Interpret types as *fibrations*

- $\Gamma \vdash A \text{ type} \rightsquigarrow$ family of cubical sets indexed by $\llbracket \Gamma \rrbracket$
that “respect paths in $\llbracket \Gamma \rrbracket$ ”

Fibrations

▣ Part 1 (coercion):

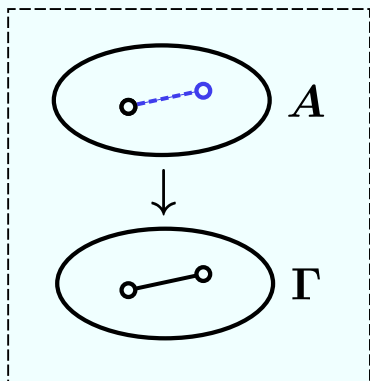
“if $\gamma_0 \text{ --- } \gamma_1 \in \llbracket \Gamma \rrbracket$ then $\llbracket A \rrbracket(\gamma_0) \simeq \llbracket A \rrbracket(\gamma_1)$ ”



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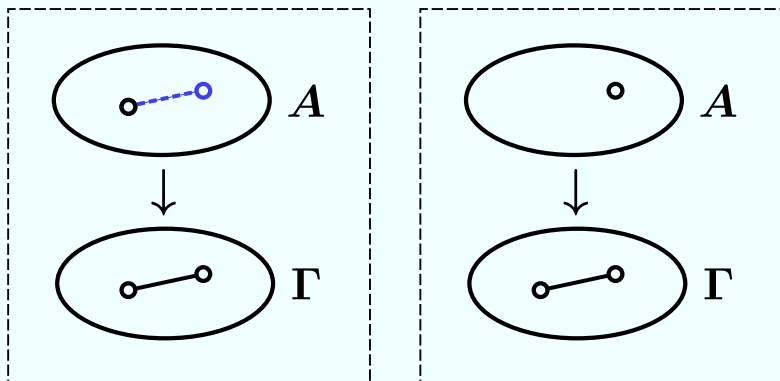
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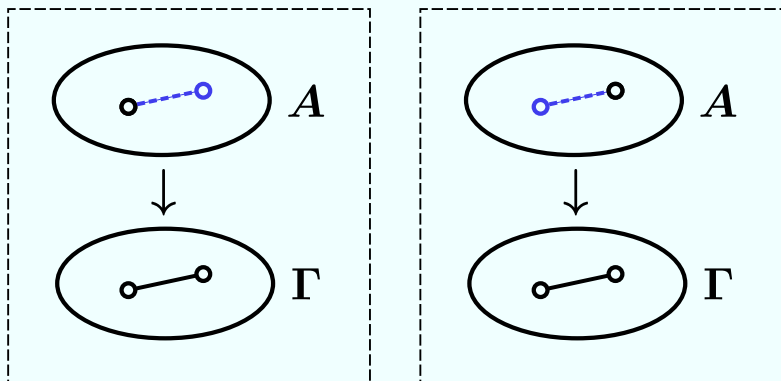
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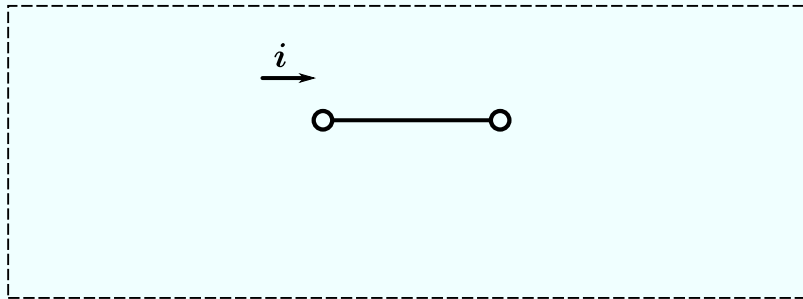
▣ Part 2 (composition): a cube in A can be adjusted

Fibrations

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“if $\gamma_0 \multimap \gamma_1 \in \llbracket \Gamma \rrbracket$ then $\llbracket A \rrbracket(\gamma_0) \simeq \llbracket A \rrbracket(\gamma_1)$ ”

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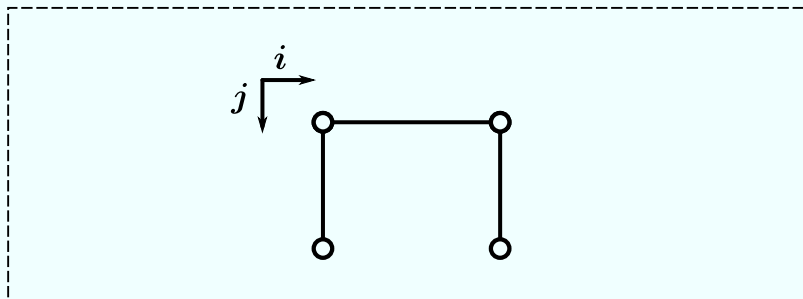


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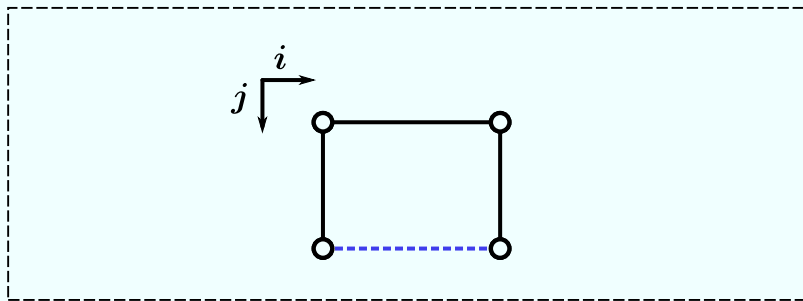


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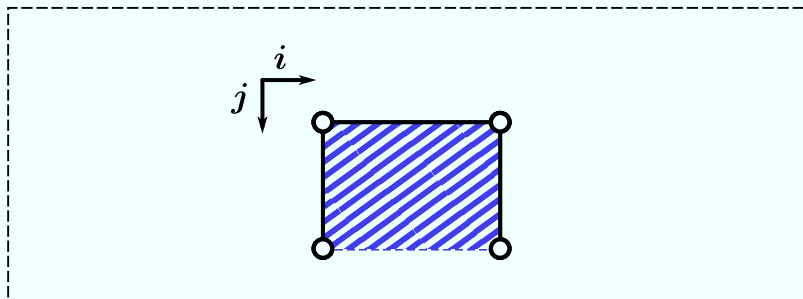


Fibrations

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Fibrations

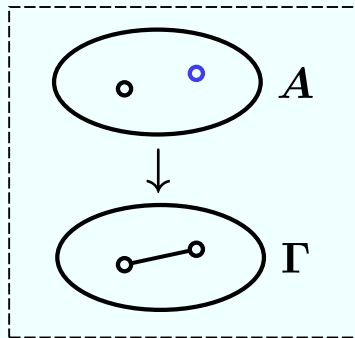
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- ▣ Part 2 (composition): a cube in A can be adjusted
- ▣ A **fibration** is a family supporting these operations

Two approaches

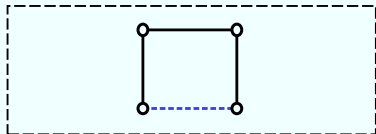
▣ Cohen, Coquand, Huber, & Mörtberg '15



coercion $0 \rightarrow 1$

$$\neg r, r \sqcup s, r \sqcap s \in \mathbb{I}$$

negation and min/max (**connection**)
operations on \mathbb{I}

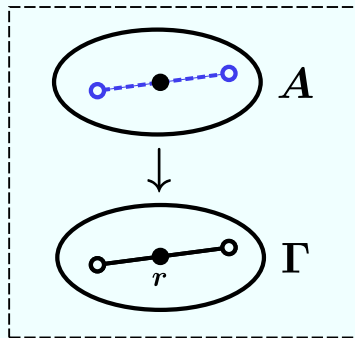


composition adjusts faces

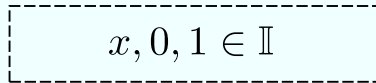
▣ Result: fibrations closed under type formers

Two approaches

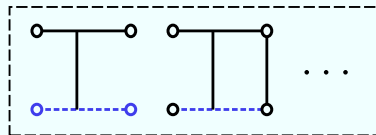
▣ Angiuli, Favonia, & Harper '18



coercion $r \rightarrow s$



only variables and faces in \mathbb{I}



composition adjusts any part

▣ Result: fibrations closed under type formers

Two approaches

▨ CCHM

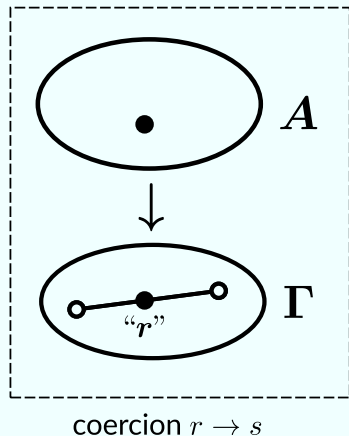
- ◆ Rely on operations on dimensions (\sqcup, \sqcap, \neg) to show closure under type formers
- ◆ Does not apply in AFH cubical sets

▨ AFH

- ◆ Need stronger composition and coercion to show closure under type formers
- ◆ Applies in CCHM cubical sets, but gives inequivalent definition of fibration

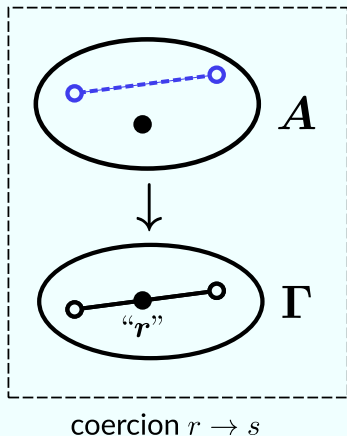
Is there a unifying construction that generalizes these?

Unifying construction



- Q:** Where do we use the stronger composition in AFH?
- A:** Fixing coercion output that doesn't quite agree with input

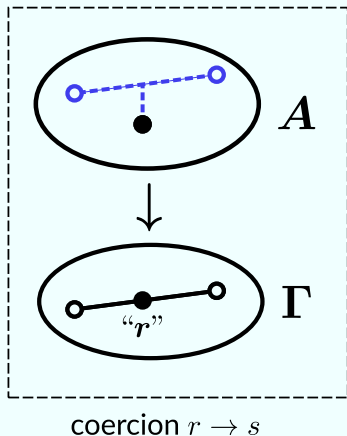
Unifying construction



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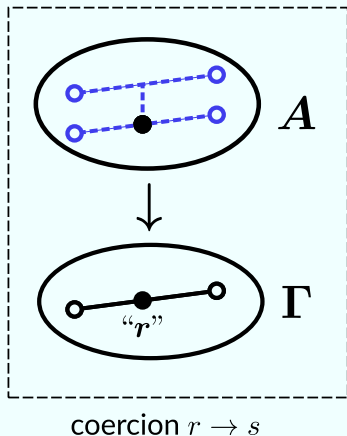
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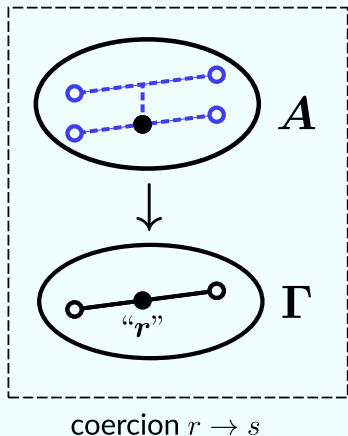
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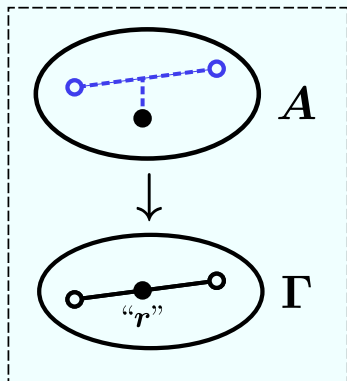
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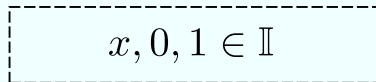
IDEA (CMS):

Weaken the condition on coercion output

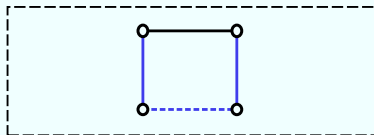
Unifying construction



weak coercion $r \rightarrow s$



only variables and faces
required in \mathbb{I}



composition adjusts faces
+ arbitrary specified Φ

- ▨ Fibrations are closed under type formers
- ▨ Fibrations participate in a model structure

Unifying construction

- Parameterized by category \mathcal{C} with \mathbb{I} and Φ (+ axioms)

$$\text{AFH} = \text{CMS} \left(\begin{array}{c} \text{cartesian} \\ \text{cubical sets} \end{array}, \mathbb{I}, \begin{array}{c} \circ \text{---} \circ \\ | \\ \circ \text{---} \circ \end{array} \right)$$

(strict and weak coercion become equivalent)

$$\text{CCHM} = \text{CMS} \left(\begin{array}{c} \text{De Morgan} \\ \text{cubical sets} \end{array}, \mathbb{I}, \begin{array}{c} \circ \text{---} \circ \\ | \quad | \\ \circ \text{---} \circ \end{array} \right)$$

(\neg, \sqcup, \sqcap + coercion $0 \rightarrow 1 \Rightarrow$ weak coercion)

- Also new models, e.g. cartesian w/ only faces in Φ

Unifying construction

- ▣ Formulated following Orton & Pitts '16 (for CCHM), Angiuli, Brunerie, Coquand, Favonia, Harper, & Licata '18 (for AFH)
 - ◆ assume \mathcal{C} interprets ordinary type theory
 - ◆ describe axioms and construction in internal language
 - ◆ enables straightforward formalization (ours in Agda)

Unifying construction

▨ Model structure

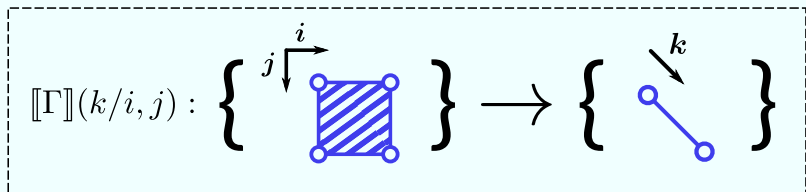
- ◆ setting for homotopy theory
 - ◆ following Sattler '17 (for CCHM)
 - ◆ use Swan '18 to translate coercion $r \rightarrow s$
- $(C, W, F) \begin{cases} C \text{ (cofibrations): generated by } \Phi \\ W \text{ (weak equivalences): equivalences} \\ F \text{ (fibrations): fibrations} \end{cases}$
- ◆ Our (C, W, F) has F maximal such that families in F have coercion $0 \rightarrow r$

Future work

▣ Original cubical model: Bezem, Coquand, & Huber '13

- ◆ Substructural: no diagonal maps between cubes

$$(k/i, j) : (k : \mathbb{I}) \rightarrow (i : \mathbb{I}, j : \mathbb{I})$$



- ◆ Definitions of fibration structure for types rely on the *absence* of diagonals

▣ How do cubical models relate to other models?